

# INTRODUCTION TO GATE-BASED QUANTUM COMPUTING

**JUNIQ Spring School on Quantum Computing 2025** 

31 MARCH 2025 | DR. DENNIS WILLSCH

WE COMPUTE WITH QUANTUM COMPUTERS



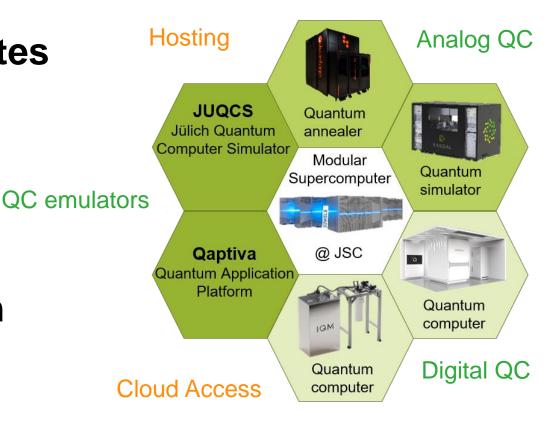






https://arxiv.org/pdf/2201.02051.pdf

- 1. Quantum Bits and Quantum Gates
- 2. Programming and Simulating Quantum Circuits
- 3. Applications:
  - 1. Quantum Fourier Transform
  - 2. Quantum Adder
  - 3. Quantum Approximate Optimization Algorithm









#### A single qubit

> Definition of a single qubit

$$|\psi\rangle = \psi_0 |0\rangle + \psi_1 |1\rangle = \begin{pmatrix} \psi_0 \\ \psi_1 \end{pmatrix}$$

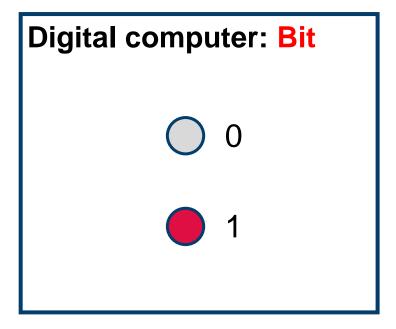
Computational basis states:

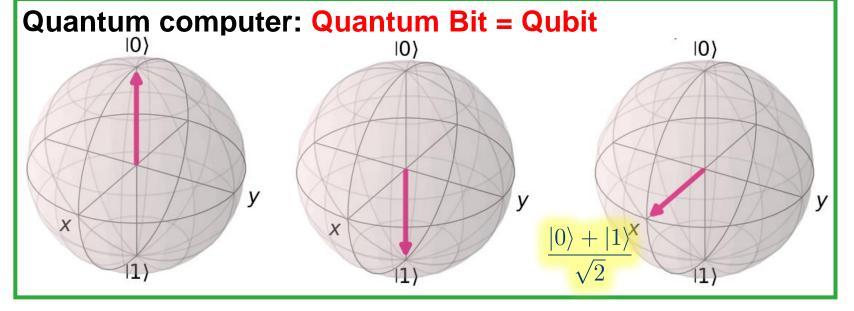
$$|0\rangle = \begin{pmatrix} 1\\0 \end{pmatrix} \qquad |1\rangle = \begin{pmatrix} 0\\1 \end{pmatrix}$$

#### Important qubits:

$$|+\rangle = \frac{1}{\sqrt{2}} |0\rangle + \frac{1}{\sqrt{2}} |1\rangle$$

$$\left|-\right\rangle = \frac{1}{\sqrt{2}} \left|0\right\rangle - \frac{1}{\sqrt{2}} \left|1\right\rangle$$





#### A single qubit

> Definition of a single qubit

$$|\psi\rangle = \psi_0 |0\rangle + \psi_1 |1\rangle = \begin{pmatrix} \psi_0 \\ \psi_1 \end{pmatrix}$$

> Representation on the Bloch sphere

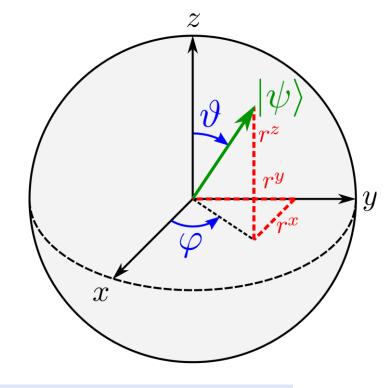
$$|\psi\rangle = \cos\frac{\vartheta}{2}|0\rangle + e^{i\varphi}\sin\frac{\vartheta}{2}|1\rangle = \begin{pmatrix} \cos\frac{\vartheta}{2} \\ e^{i\varphi}\sin\frac{\vartheta}{2} \end{pmatrix}$$

#### Normalization:

$$\langle \psi | \psi \rangle = |\psi_0|^2 + |\psi_1|^2 = 1$$
  
 $\langle \psi | = (\psi_0^*, \ \psi_1^*)$ 

#### Global phase:

$$|\psi\rangle \equiv e^{i\Phi}|\psi\rangle$$



> Evaluation of the Cartesian coordinates

$$\vec{r} = \begin{pmatrix} r^x \\ r^y \\ r^z \end{pmatrix} = \begin{pmatrix} \langle \psi | \sigma^x | \psi \rangle \\ \langle \psi | \sigma^y | \psi \rangle \\ \langle \psi | \sigma^z | \psi \rangle \end{pmatrix} = \begin{pmatrix} \sin \vartheta \cos \varphi \\ \sin \vartheta \sin \varphi \\ \cos \vartheta \end{pmatrix}$$

#### Pauli matrices:

$$\sigma^x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \sigma^y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \quad \sigma^z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$



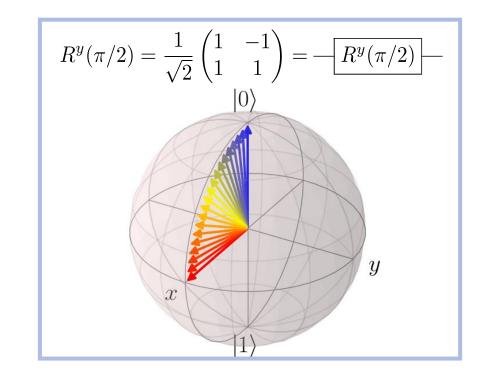
#### Gates as rotations on the Bloch sphere

> Rotations around the axes of the Bloch sphere

$$R^{x}(\theta) = e^{-i\theta\sigma^{x}/2} = \begin{pmatrix} \cos(\theta/2) & -i\sin(\theta/2) \\ -i\sin(\theta/2) & \cos(\theta/2) \end{pmatrix}$$

$$R^{y}(\theta) = e^{-i\theta\sigma^{y}/2} = \begin{pmatrix} \cos(\theta/2) & -\sin(\theta/2) \\ \sin(\theta/2) & \cos(\theta/2) \end{pmatrix}$$

$$R^{z}(\theta) = e^{-i\theta\sigma^{z}/2} = \begin{pmatrix} \exp(-i\theta/2) & 0\\ 0 & \exp(i\theta/2) \end{pmatrix}$$



#### General rotation gate

$$R^{\vec{n}}(\theta) = e^{-i\theta\vec{n}\cdot\vec{\sigma}/2} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}\vec{n}\cdot\vec{\sigma}$$

#### Standard gate set:

$$X = \sigma^x Y = \sigma^y Z = \sigma^z$$

$$Y = \sigma^{y}$$

$$Z = \sigma^z$$

$$H=\frac{1}{\sqrt{2}}\left($$

$$S = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \qquad S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix} \qquad T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$$

#### Multiple qubits

Multi-qubit state

$$|\psi\rangle = \sum_{j=0}^{2^n-1} \psi_j \, |j\rangle = \psi_0 \, |0\cdots 00\rangle + \psi_1 \, |0\cdots 01\rangle + \cdots + \psi_{2^n-1} \, |1\cdots 11\rangle = \begin{pmatrix} \psi_0 \\ \psi_1 \\ \vdots \\ \psi_{2^n-1} \end{pmatrix}$$

$$|\phi\rangle = \sum_{j=0}^{2^n-1} \psi_j \, |j\rangle = \psi_0 \, |0\cdots 00\rangle + \psi_1 \, |0\cdots 01\rangle + \cdots + \psi_{2^n-1} \, |1\cdots 11\rangle = \begin{pmatrix} \psi_0 \\ \psi_1 \\ \vdots \\ \psi_{2^n-1} \end{pmatrix}$$

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$$|\phi\rangle = \begin{pmatrix} \psi_0 \\ \psi_1 \\ \vdots \\ \psi_{2^n-1} \end{pmatrix}$$

$$|\phi\rangle = \langle \psi\rangle \otimes |\phi\rangle \otimes |\phi\rangle$$

➤ Important two-qubit gates

$$CNOT = CX = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix} = \begin{bmatrix} U \\ U \\ U \end{bmatrix} = \begin{bmatrix} I & 0 \\ 0 & U \\ U \end{bmatrix} = \begin{bmatrix} I & 0 \\$$

$$CZ = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} =$$

General controlled gates

$$CU_{i_{1}i_{2}} | q_{0} \cdots q_{n-1} \rangle = \begin{cases} |q_{0} \cdots q_{n-1} \rangle & \text{(if } q_{i_{1}} = 0) \\ |q_{0} \cdots q_{i_{2}-1} \rangle \left( U | q_{i_{2}} \rangle \right) | q_{i_{2}+1} \cdots q_{n-1} \rangle & \text{(if } q_{i_{1}} = 1) \end{cases}$$

#### Computational basis states:

$$|00\rangle = \begin{pmatrix} 1\\0\\0\\0 \end{pmatrix} \quad |01\rangle = \begin{pmatrix} 0\\1\\0\\0 \end{pmatrix} \quad |10\rangle = \begin{pmatrix} 0\\0\\1\\0 \end{pmatrix} \quad |11\rangle = \begin{pmatrix} 0\\0\\0\\1 \end{pmatrix}$$

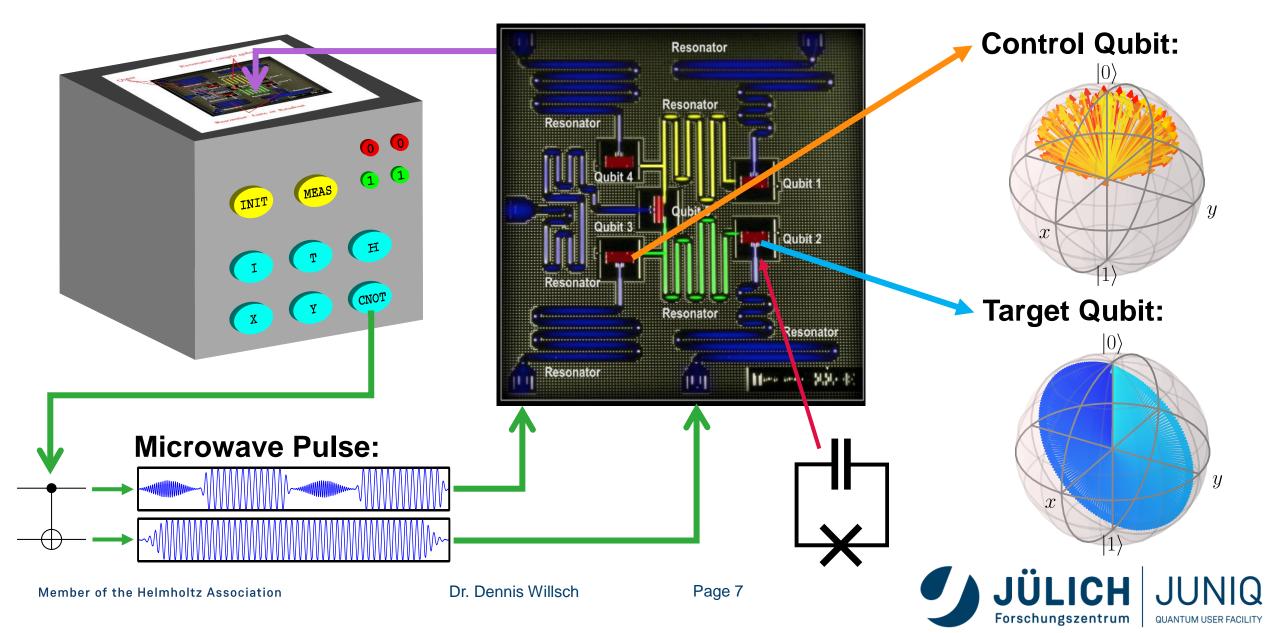
$$|q_0q_1\rangle = |q_0\rangle \otimes |q_1\rangle$$
  $A \otimes B = \begin{pmatrix} a_{00}B & a_{01}B & \cdots \\ a_{10}B & a_{11}B & \\ \vdots & & \ddots \end{pmatrix}$ 

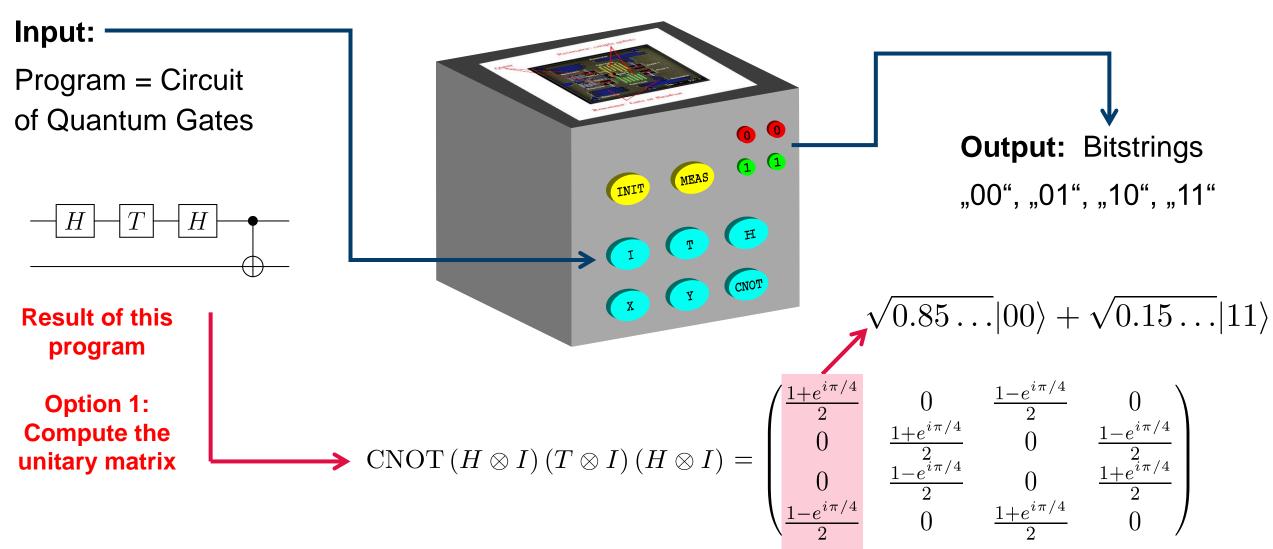
$$\mathbf{C}\mathbf{U} = \begin{pmatrix} I & 0 \\ 0 & \mathbf{U} \end{pmatrix} = \mathbf{U}$$

$$SWAP = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} =$$

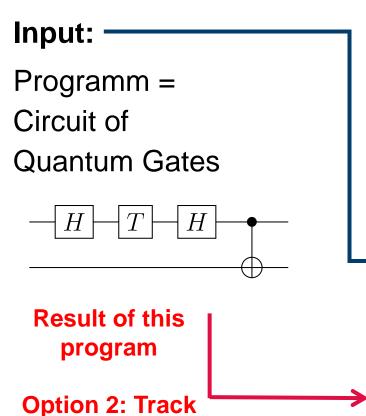


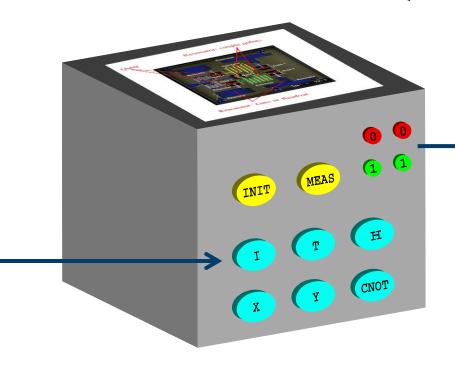
# INTERLUDE: SUPERCONDUCTING QUANTUM COMPUTERS











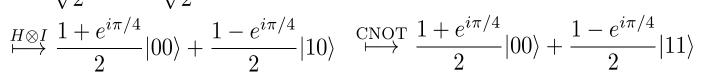
**Output:** Bitstrings

"00", "01", "10", "11"

 $\sqrt{0.85\ldots|00\rangle} + \sqrt{0.15\ldots}|11\rangle$ 

 $|00\rangle \stackrel{H\otimes I}{\longmapsto} \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|10\rangle$  $\stackrel{T\otimes I}{\longrightarrow} \frac{1}{\sqrt{2}}|00\rangle + \frac{e^{i\pi/4}}{\sqrt{2}}|10\rangle$ 

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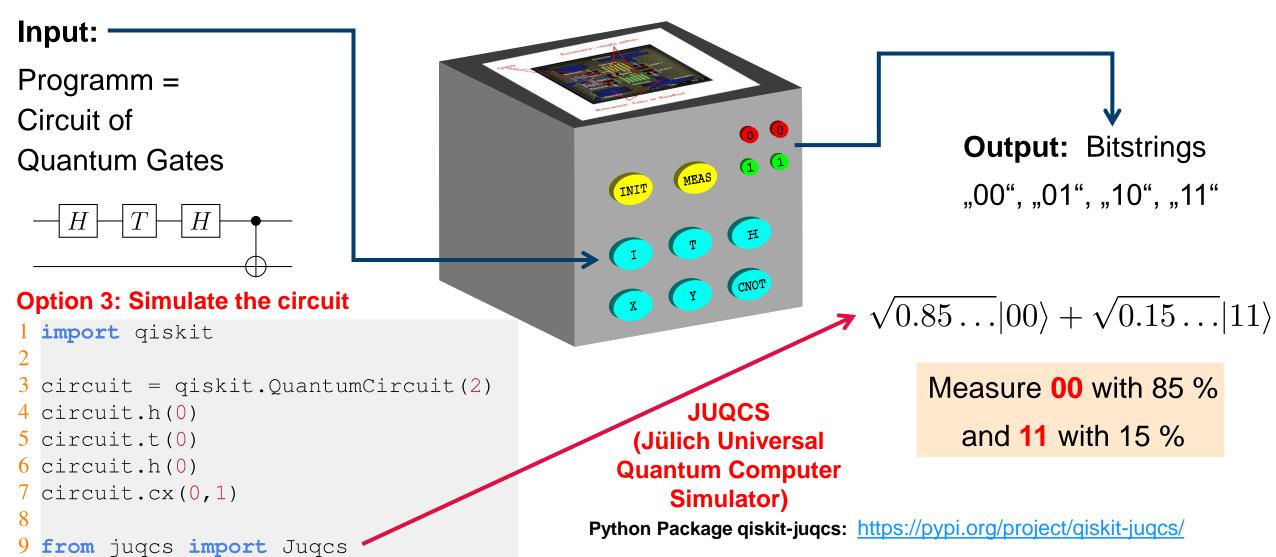






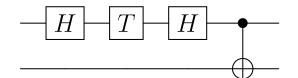
the unitary

transformations



#### Simulation with Qiskit Aer

qasm\_simulator



statevector\_simulator

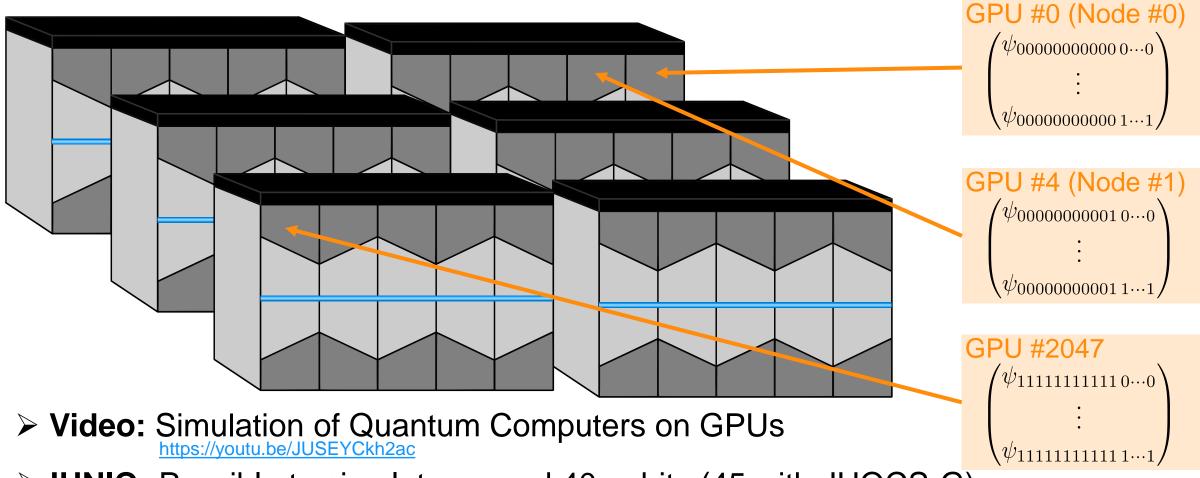
```
import qiskit
circuit = qiskit.QuantumCircuit(2)
circuit.h(0)
circuit.t(0)
circuit.h(0)
circuit.cx(0,1)
circuit.measure all()
backend = qiskit.Aer.get backend('qasm simulator')
result = backend.run(circuit,
                    shots=1000).result()
result.get counts()
{'00': 846, '11': 154}
                  Measure 00 with 85 %
                     and 11 with 15 %
```

```
import qiskit
circuit = qiskit.QuantumCircuit(2)
circuit.h(0)
circuit.t(0)
circuit.h(0)
circuit.cx(0,1)
backend = qiskit.Aer.get backend('statevector simulator')
result = backend.run(circuit.reverse bits()).result()
result.get statevector()
Statevector([0.85355339+0.35355339j, 0.
                                               +0.j
                       +0.j
                                   , 0.14644661-0.35355339i],
             0.
            \sqrt{0.85...|00} + \sqrt{0.15...|11}
```



## SIMULATING QUANTUM CIRCUITS

Simulation with JUQCS (large-scale simulations on a supercomputer)



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>JUNIQ: Possible to simulate around 40 qubits (45 with JUQCS-G)

https://juniq.fz-juelich.de

# **QUANTUM FOURIER TRANSFORM**

An operation to move registers into phases and back

> The formal definition of the QFT is:

Here happens a multiplication modulo 2<sup>n</sup>

→ Can (ab)use QFT for arithmetics! (addition, multiplication, etc.)

$$|j_0 j_1 j_2 \cdots\rangle \stackrel{\text{QFT}}{\mapsto} \frac{1}{\sqrt{2^n}} \sum_{k_0 k_1 k_2 \cdots} e^{2\pi i (j_0 j_1 j_2 \cdots) \cdot (k_0 k_1 k_2 \cdots)/2^n} |k_0 k_1 k_2 \cdots\rangle$$

> Intuition:

$$|j\rangle \leftrightarrow \sum_{k=0}^{2^n-1} e^{2\pi i j k/2^n} |k\rangle$$

> Implementation for two qubits:

$$|j_0\rangle$$
  $H$   $S$   $H$   $H$ 

Verification:

$$\text{SWAP } (I \otimes H) \text{ CS } (H \otimes I) \ket{j_0 j_1} = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix} \ket{j_0 j_1} = \frac{1}{2} \sum_{k_0 k_1} e^{2\pi i \, (j_0 j_1) \cdot (k_0 k_1)/4} \ket{k_0 k_1}$$

$$\frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix}$$

$$e^{2\pi i (j_0 j_1) \cdot 0/4} |0\rangle$$

$$+ e^{2\pi i (j_0 j_1) \cdot 1/4} |1\rangle$$

$$+ e^{2\pi i (j_0 j_1) \cdot 2/4} |2\rangle$$

$$+ e^{2\pi i (j_0 j_1) \cdot 3/4} |3\rangle$$

$$|j_0 j_1\rangle = \frac{1}{2} \sum_{k_0 k_1} e^{2\pi i (j_0 j_1) \cdot (k_0 k_1)/4} |k_0 k_1\rangle$$





## **QUANTUM ADDER**

#### A nice application of the QFT

We are looking for a quantum circuit to implement the following unitary operation

$$|x_0x_1\rangle|y_0y_1\rangle$$
 $\mapsto$ 
 $|x_0x_1\rangle|y_0y_1+x_0x_1\rangle$ 
 $\xrightarrow{\text{Sum } s_0s_1}$ 

- > Such a modulo-4 adder would also work on superpositions!
- > Examples:

$$|2\rangle |1\rangle \qquad \mapsto \qquad |2\rangle |3\rangle$$

$$|2\rangle \frac{|0\rangle + |1\rangle}{\sqrt{2}} \qquad \mapsto \qquad |2\rangle \frac{|2\rangle + |3\rangle}{\sqrt{2}}$$

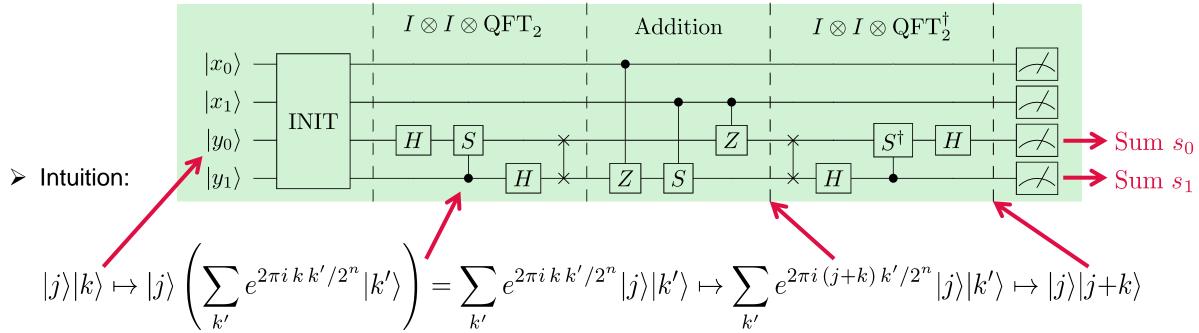
$$|2\rangle \frac{|0\rangle + |1\rangle + |2\rangle}{\sqrt{3}} \qquad \mapsto \qquad |2\rangle \frac{|2\rangle + |3\rangle + |0\rangle}{\sqrt{3}}$$

$$|0\rangle + |1\rangle \frac{|2\rangle + |3\rangle + |0\rangle}{\sqrt{2}} \qquad \mapsto \qquad \frac{|0\rangle |2\rangle + |0\rangle |3\rangle + |0\rangle |0\rangle + |1\rangle |3\rangle + |1\rangle |0\rangle + |1\rangle |1\rangle}{\sqrt{6}}$$

### **QUANTUM ADDER**

#### A nice application of the QFT

- ➤ Idea: → The QFT converts between registers and phases
  - → By using bitwise phase shifts, we can implement the addition in the exponent
  - $\rightarrow$  This is automatically modulo 4 ( $\leftarrow$  periodicity of the complex exponential function)
- > A circuit to realize this idea is:



## **QUANTUM APPROXIMATE OPTIMIZATION ALGORITHM**

#### A very brief overview of the QAOA

Say we want to find the minimum of the function

$$E(q_0,\ldots,q_{n-1})=\sum_{i,j}q_iQ_{ij}q_j \qquad \Leftrightarrow \qquad E(s_0,\ldots,s_{n-1})=\sum_{i< j}h_is_i+\sum_{i< j}J_{ij}s_is_j$$
 "QUBO" (q = 0,1)

- Combinatorial optimization problem
  - → discrete optimization is hard!
- ➤ Idea: With a quantum circuit, we can create a superposition

$$|\psi\rangle = \psi_0 |0\cdots 00\rangle + \cdots + \psi_{q_0^*\cdots q_{n-1}^*} |q_0^*\cdots q_{n-1}^*\rangle + \cdots + \psi_{2^n-1} |1\cdots 11\rangle$$

Solution to the problem

Goal: Enhance this term!



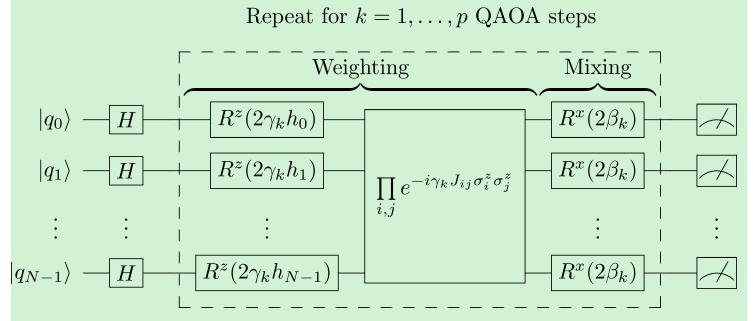
### **QUANTUM APPROXIMATE OPTIMIZATION ALGORITHM**

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- $\triangleright$  Make the quantum circuit dependent on real parameters  $(\beta_1,\ldots,\beta_p,\gamma_1,\ldots,\gamma_p)$
- > Hope: these 2p parameters are easier searchable than the original discrete variables



# **QUANTUM APPROXIMATE OPTIMIZATION ALGORITHM**

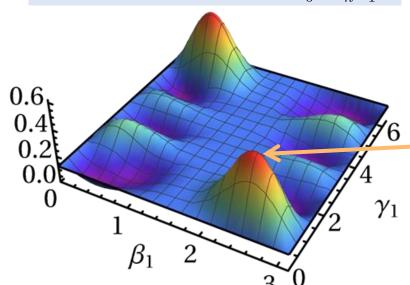
#### A very brief overview of the QAOA

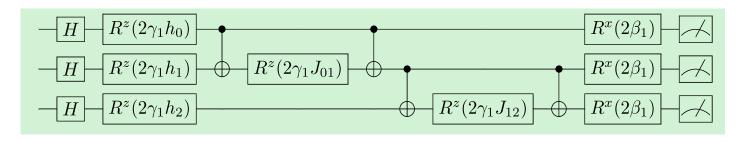
Goal: Enhance this term!

$$|\psi\rangle = \psi_0 |0\cdots 00\rangle + \cdots + \psi_{q_0^*\cdots q_{n-1}^*} |q_0^*\cdots q_{n-1}^*\rangle + \cdots + \psi_{2^n-1} |1\cdots 11\rangle$$

➤ Look at p=1 QAOA step

Success probability  $|\psi_{q_0^*\cdots q_{n-1}^*}|^2$ 





There are regions where the success probability is enhanced! How to find them?

> Will be a dedicated topic in this School ©



# THANK YOU FOR YOUR ATTENTION

> Further information:

Programming Quantum Computers: <a href="https://arxiv.org/pdf/2201.02051.pdf">https://arxiv.org/pdf/2201.02051.pdf</a>

> JUQCS: Video <a href="https://youtu.be/JUSEYCkh2ac">https://youtu.be/JUSEYCkh2ac</a>

+ **JUQCS:** Paper <a href="https://doi.org/10.1016/j.cpc.2022.108411">https://doi.org/10.1016/j.cpc.2022.108411</a>